U.S. Pat. App. Ser. N(---)/090,718 Att. Docket No. 10191/2275 Reply to Final Office Action of 05/10/2007

AMENDMENTS to the CLAIMS:

Without prejudice, this listing of the claims replaces all prior versions and listings of the claims in the present application:

LISTING OF CLAIMS:

1. (Currently Amended) A method of data encryption in programming of a control unit comprising:

encrypting a complete stream of data to be transmitted in a programming unit using a first key, wherein a byte by byte encryption of the complete stream of data is capable of being performed, and wherein no byte-wise allocation between input and output data occurs;

transmitting the data that had been encrypted to the control unit via a data line; and decrypting the data that had been encrypted in the programming unit using a second key provided in the control unit;

wherein:

successive bytes during encryption are provided with an index i, where i = 0, $1, 2, \ldots$,

an encrypted byte n^* is formed from an unencrypted byte n according to the following, a starting value n_{-1} being used for decryption and encryption:

$$\frac{n_{-1} \equiv S_o}{n_i^* = \left(n_i < < \sum_{j=0}^{i} n_{j-1}^*\right) \oplus S_{h\left(\sum_{j=0}^{i} n_{j-1}^*\right)}}$$

an unencrypted byte n is formed from an encrypted byte n* according to the following:

$$n_i = \left(n_i^* \oplus S_{h\left(\sum\limits_{j=0}^i n_{j-1}^*\right)}\right) >>> \sum\limits_{j=0}^i n_{j-1}^*$$

2. (Original) The method of claim 1, wherein the first key and the second key are identical.

U.S. Pat. App. Ser. N(-)/090,718 Att. Docket No. 10191/2275 Reply to Final Office Action of 05/10/2007

- 3. (Original) The method of claim 1, wherein the first key and the second key are not identical.
- 4. (Original) The method of claim 2, wherein each one of the first key and the second key includes a table that is accessed by a hash function.
- 5. (Original) The method of claim 1, wherein at least one of the first key and the second key is implemented in an electronic circuit.
- 6. (Original) The method of claim 1, wherein at least one of the first key and the second key is implemented in the form of a computer program.
- 7. (Currently Amended) A data encryption system, comprising:
 - a programming unit in which a first key is provided;
 - a control unit in which a second key is provided; and
- a data line coupled to the programming unit and the control unit for transmitting encrypted data, the encrypted data being an encryption of a complete stream of data, wherein a byte by byte encryption of the complete stream of data is capable of being performed, wherein encryption of a byte includes a rotation of bits of the byte about a number of positions, the number depending on an entire history of the encryption of the data, and wherein no byte-wise allocation between input and output data occurs;

wherein:

successive bytes during encryption are provided with an index i, where i = 0, $1, 2, \ldots$,

an encrypted byte n^* is formed from an unencrypted byte n according to the following, a starting value n_{-1} being used for decryption and encryption:

$$\frac{n_{-1} \equiv S_o}{n_i^* = \left(n_i < << \sum_{j=0}^{i} n_{j-1}^*\right) \oplus S_{h\left(\sum_{j=0}^{i} n_{j-1}^*\right)}}$$

U.S. Pat. App. Ser. No. 1090,718 Att. Docket No. 10191/2275 Reply to Final Office Action of 05/10/2007

an unencrypted byte n is formed from an encrypted byte n* according to the following:

$$n_{i} = \left(n_{i}^{*} \oplus S_{h\left(\sum\limits_{j=0}^{i} n_{j-1}^{*}\right)}\right) >>> \sum\limits_{j=0}^{i} n_{j-1}^{*}$$

- 8. (Original) The system of claim 7, wherein the first key and the second key are identical.
- 9. (Original) The system of claim 7, wherein the first key and the second key are not identical.
- 10. (Original) The system of claim 7, wherein the programming unit and the control unit each includes an electronic computing unit and a memory module that are linked together by a data bus.
- 11. (Previously Presented) A computer program product having program code executable by a computing unit, the program code when executed causing the computing unit to perform a method, the method comprising:

performing an encryption of a complete stream of data in accordance with a table and a hash function, wherein a byte by byte encryption of the complete stream of data is capable of being performed, and wherein no byte-wise allocation between input and output data occurs.

12. (Previously Presented) The computer program product of claim 11, wherein the computing unit includes an electronic computing unit in a programming unit.

U.S. Pat. App. Ser. ()0/090,718 Att. Docket No. 10191/2275 Reply to Final Office Action of 05/10/2007

13. (Previously Presented) A computer program product having program code executable by a computing unit, the program code when executed causing the computing unit to perform a method, the method comprising:

performing a decryption of a complete stream of data in accordance with a table and a hash function, wherein a byte by byte decryption of the complete stream of data is capable of being performed, and wherein no byte-wise allocation between input and output data occurs.

- 14. (Previously Presented) The computer program product of claim 13, wherein the computing unit includes an electronic computing unit in a control unit.
- 15. (Previously Presented) A computer-readable medium, comprising:

a program code executable on a computing unit for performing an encryption of a complete stream of data in accordance with a table and a hash function, wherein a byte by byte encryption of the complete stream of data is capable of being performed, and wherein no byte-wise allocation between input and output data occurs.

16. (Previously Presented) A computer-readable medium, comprising:

a program code executable on a computing unit for performing a decryption of a complete stream of data in accordance with a table and a hash function, wherein a byte by byte decryption of the complete stream of data is capable of being performed, and wherein no byte-wise allocation between input and output data occurs.

- 17. (Previously Presented) The method of claim 1, wherein there is no bit-wise allocation between input and output data.
- 18. (Previously Presented) The method of claim 7, wherein there is no bit-wise allocation between input and output data.
- 19. (Previously Presented) The method of claim 11, wherein there is no bit-wise allocation between input and output data.